

FUNDAMENTALS OF LOGIC

NO.7 PREDICATE LOGIC

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lecture URL

<https://vu5.sfc.keio.ac.jp/slide/>

So Far

- Propositional Logic
 - Logical Connectives (\wedge , \vee , \rightarrow , \neg)
 - Truth Table
 - Tautology
 - Normal Form
 - Axiom and Proof
 - LK Frame Work (Sequent Calculus)
 - Soundness and Completeness

Limitation of Propositional Logic

- Propositional Logic
 - Each proposition is either true or false.
 - The truth value does not change.
 - The truth value does not depend of objects which are referred in the proposition.
- Socrates problem:
 - Socrates is a man.
 - All men are mortal.
 - Therefore, Socrates is mortal.
- In propositional logic:
 - p = "Socrates is a man"
 - q = "All men are mortal"
 - r = "Socrates is mortal"
 - $p \wedge q \rightarrow r$?

Propositional Logic to Predicate Logic

- Extend logic to handle objects and express properties and relations of objects.
- Set of objects
 - Integer
 - Human
- Variable over a set of objects
 - **object variable**
 - x, y, z, \dots
- Name of object
 - **object constant**
 - Socrates, Pythagoras, 123, SFC, Keio, ...

Predicate

- **Predicate**

- Object x has property $P : P(x)$
- Relation R holds between object x and object $y : R(x, y)$

- $Q(x_1, x_2, \dots, x_n)$

- Q holds for objects x_1, x_2, \dots, x_n .
- Q is a predicate with n variables.

- $P(x) = "x \text{ is a man}"$

- $P(\text{Socrates}) = "Socrates \text{ is a man}"$
- $P(\text{Pythagoras}) = "Pythagoras \text{ is a man}"$
- $P(\text{Taro}) = "Taro \text{ is a man}"$

- $R(x, y) = "x \text{ likes } y"$

- $R(\text{Taro}, \text{Hanako}) = "Taro \text{ likes Hanako}"$
- $R(\text{Taro}, \text{Momoko}) = "Taro \text{ likes Momoko}"$
- $R(\text{Hanako}, \text{Taro}) = "Hanako \text{ likes Taro}"$

Quantifier

- $P(x)$
 - Which x makes P hold?
 - Does it hold for any x ?
 - Does it hold for some x ?
- **Quantifier**
 - $\forall x P(x)$
 - **Universal quantifier**
 - For any x , $P(x)$ holds.
 - $P(x)$ holds for all x .
 - $\exists x P(x)$
 - **Existential quantifier**
 - For some x , $P(x)$ holds.
 - There exists x which makes $P(x)$ hold.
- $Q(x) = "x \text{ is mortal}"$
 - $\forall x Q(x) = "Everybody \text{ is mortal}"$
 - $\exists x Q(x) = "Someone \text{ is mortal}"$, "There is someone who is mortal"

Predicate Logic

- Predicate Logic
 - Use predicates instead of propositional variables.
 - Four logical connectives: \wedge , \vee , \rightarrow , \neg
 - Two quantifiers: $\forall x$, $\exists x$
- Socrates example: $P(x)$ = "x is a man", $Q(x)$ = "x is mortal"
 - $P(\text{Socrates})$ = "Socrates is a man"
 - $\forall x(P(x) \rightarrow Q(x))$ = "All men are mortal"
 - $Q(\text{Socrates})$ = "Socrates is mortal"
- Math example: $P(x)$ = "x is a prime number bigger than 2", $Q(x)$ = "x is an odd number"
 - $P(7)$ = "7 is a prime number bigger than 2"
 - $\forall x(P(x) \rightarrow Q(x))$ = "Any prime number bigger than 2 is an odd number"
 - $Q(7)$ = "7 is an odd number"

Example (1)

- Let $S(x)$ and $M(x)$ be as follows:

- $S(x)$ = "x is an SFC student"
- $M(x)$ = "x likes math"

- Write the meaning of the following formulae:

- $\forall x S(x)$ = " "
- $\exists x S(x)$ = " "
- $\forall x(S(x) \rightarrow M(x))$ = " "
- $\exists x(S(x) \wedge M(x))$ = " "
- $\forall x(S(x) \rightarrow \neg M(x))$ = " "
- $\forall x(\neg S(x) \rightarrow M(x))$ = " "
- $\forall x \neg S(x)$ = " "
- $\neg \forall x S(x)$ = " "
- $\neg \forall x(S(x) \rightarrow M(x))$ = " "
- $\forall x \neg(S(x) \rightarrow M(x))$ = " "
- $\exists x \neg S(x)$ = " "
- $\neg \exists x S(x)$ = " "

Example (2)

- Let $L(x, y)$ mean " x likes y ". Write the meaning of the following formulae:

- $\forall x L(\text{Taro}, x)$ = "
- $\exists x L(\text{Taro}, x)$ = "
- $\forall x L(x, \text{Taro})$ = "
- $\exists x L(x, \text{Taro})$ = "
- $\forall x \forall y L(x, y)$ = "
- $\exists x \exists y L(x, y)$ = "
- $\forall x \exists y L(x, y)$ = "
- $\exists x \forall y L(x, y)$ = "
- $\exists y \forall x L(x, y)$ = "
- $\forall x \forall y (S(x) \rightarrow L(x, y))$ = "
- $\forall x \forall y (S(y) \rightarrow L(x, y))$ = "
- $\forall x (S(x) \rightarrow \forall y L(x, y))$ = "
- $\forall x (\forall y L(x, y) \rightarrow S(x))$ = "

Language for Predicate Logic

- A set of symbols for predicate logic is called **language**.
 - It is different from linguistic language.
 - It is closer to **vocabulary**.
- A language L of predicate logic consists of the followings:
 1. Logical connectives: $\wedge, \vee, \rightarrow, \neg$
 2. Quantifiers: \forall, \exists
 3. Object variables: x, y, z, \dots
 4. Object constants: c, d, \dots
 5. Function symbols: f, g, \dots
 6. Predicate symbols: P, Q, \dots

Terms

- **Terms** of a language L is defined as follows:
 1. Object variables and constants of L are terms.
 2. For a function symbol f and m variables (arity m) in L , if t_1, \dots, t_m are terms, $f(t_1, \dots, t_m)$ is also a term.
- **Example: Natural Number Theory**
 - Object constants: 0, 1, etc.
 - Function symbols: $S(x)$, $+$, \times , etc.
 - Predicate symbols: $=$, $<$, etc.
 - Terms
 - x
 - 0
 - $S(x) + (1 \times S(S(y)))$

Logical Formulae

- **Logical Formulae** of L are defined as follows:
 1. For a predicate symbol P of n variables in L , if t_1, \dots, t_n are terms, $P(t_1, \dots, t_n)$ is a formula (**atomic formula**).
 2. For formulae A and B , $(A \wedge B)$, $(A \vee B)$, $(A \rightarrow B)$ and $(\neg A)$ are formulae.
 3. For a formula A and an object variable x , $(\forall x A)$ and $(\exists x A)$ are formulae.
- **Example: Natural Number Theory**
 - $\exists x(x \times z = y)$
 - $\forall x \forall y((x + S(y)) = S(x + y))$

Bound and Free Variables

- **Bound variable**
 - In $\exists z(x \times z = y)$, z of $x \times z = y$ is **bound** by $\exists z$.
 - Bound variables can be renamed without changing the meaning.
 - $\exists w(x \times w = y)$
- Variables which are not bound are **free variables**.
 - In $\exists z(x \times z = y)$, x and y are free variables.
- Variables may be bound or free depending on their **occurrence**.
 - $\exists z(x \times z = y) \wedge \exists y(x + x = y)$

Closed Formulae

- When a logical formula A does not contain free variables, A is called a **closed** logical formula.
 - $\forall x(S(x) \rightarrow \forall y L(x, y))$
- If x_1, \dots, x_n are free variables of a logical formula A ,
 - $\forall x_1 \dots \forall x_n A$
 - is called **universal closure** of A .
- In mathematics, universal quantifiers are often omitted.
 - Commutative law of addition $x + y = y + x$
 - Its universal closure: $\forall x \forall y (x + y = y + x)$

Assignment of Terms

- For a logical formula A , when all the free occurrence of x are replaced with a term t , it is called an **assignment** of t to x .
 - $A[t/x]$
- Example
 - Let A be $\exists z(x \times z = y)$.
 - $A[w/y]$ is $\exists z(x \times z = w)$.
 - $A[x/y]$ is $\exists z(x \times z = x)$.
 - $A[(x + w)/x]$ is $\exists z((x + w) \times z = y)$.
- If bound relationship is affected by an assignment, the bound variable must be changed before the assignment.
 - $A[z/y]$ is not $\exists z(x \times z = z)$, but $\exists w(x \times w = z)$.
 - In general, $(\forall xA)[t/x]$ is $\forall u(A[u/x][t/x])$ where u is a variable which does not occur in A or t .

Sub-formulae

- Define **sub-formulae** similar to propositional logic.
 1. A is a sub-formula of A .
 2. A and B are sub-formulae of $(A \wedge B)$.
 3. A and B are sub-formulae of $(A \vee B)$.
 4. A and B are sub-formulae of $(A \rightarrow B)$.
 5. A is a sub-formula of $(\neg A)$.
 6. For any term t , $A[t/x]$ is a sub-formula of $\forall xA$.
 7. For any term t , $A[t/x]$ is a sub-formula of $\exists xA$.
- When a formula contains quantifiers, there are in finitely many sub-formulae.
 - Sub-formulae of $\forall xQ(x)$ are:
 $\forall xQ(x)$, $Q(\text{Socrates})$, $Q(\text{Taro})$, $Q(\text{mother}(\text{Taro}))$, ...

Summary

- Predicate Logic
 - Limitation of propositional logic
 - Description about objects
- Logical Formulae of Predicate Logic
 - Language
 - Term
 - Logical Formulae
- Quantifiers
 - Bound and free variables
 - Closed formulae
 - Universal closure
 - Assingment