Optimization Theory (DS2) HW#2 Certificates, Standard Equality Form and Simplex Iteration

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Problems and page numbers from the paperback edition of *A Gentle Introduction* to *Optimization*, by B. Guenin *et al.*

- 1. (Problem 1 in Sec. 2.1 of the text, p. 50.)
 - (a) Prove that the following LP problem is infeasible:

$$\max\{3x_1 + 4x_2 + 6x_3\}\tag{1}$$

subject to

$$3x_1 + 5x_2 - 6x_3 = 4 \tag{2}$$

$$x_1 + 3x_2 - 4x_3 = 2 \tag{3}$$

$$-x_1 + x_2 - x_3 = -1 \tag{4}$$

$$x_1, x_2, x_3 \ge 0.$$
 (5)

(b) Prove that the following LP problem is unbounded:

$$\max\{-x_3 + x_4\}\tag{6}$$

subject to

$$x_1 + x_3 - x_4 = 1 \tag{7}$$

$$x_2 + 2x_3 - x_4 = 2 \tag{8}$$

- $x_1, x_2, x_3, x_4 \ge 0. \tag{9}$
- (c) **Optional:** Prove that the LP Problem

$$max\{\vec{c}^{\mathsf{T}}\vec{x}: A\vec{x} = \vec{b}, \vec{x} \ge \vec{0}\}$$
(10)

is unbounded, where

$$A = \begin{pmatrix} 4 & 2 & 1 & -6 & -1 \\ -1 & 1 & -4 & 1 & 3 \\ 3 & -6 & 5 & 3 & -5 \end{pmatrix},$$
 (11)

$$\vec{b} = \begin{pmatrix} 11\\ -2\\ -8 \end{pmatrix}, \tag{12}$$

$$\vec{c} = \begin{pmatrix} 1 \\ -2 \\ 1 \\ 1 \\ 1 \end{pmatrix} \tag{13}$$

Hint: Consider the vectors

$$\vec{x} = (1, 3, 1, 0, 0)^{\mathsf{T}} \text{ and } \vec{d} = (1, 1, 1, 1)^{\mathsf{T}}$$
 (14)

- (d) **Optional:** For each of the problems in parts (b) and (c), give a feasible solution having object value exactly 5000.
- 2. (Problem 1 in Sec. 2.2 of the text, p. 54.)
 - (a) Convert the following LPs into SEF:

$$\min\{(2, -1, 4, 2, 4)(x_1, x_2, x_3, x_4, x_5)^{\mathsf{T}}\}$$
(15)

subject to

$$\begin{pmatrix} 1 & 2 & 4 & 7 & 3 \\ 2 & 8 & 9 & 0 & 0 \\ 1 & 1 & 0 & 2 & 6 \\ -3 & 4 & 3 & 1 & -1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{pmatrix} \stackrel{\leq}{=} \begin{pmatrix} 1 \\ 2 \\ 3 \\ 4 \end{pmatrix}.$$
(16)
$$x_1, x_2, x_4 \ge 0.$$
(17)

(Careful, note that not all of the x_i have a non-negativity constraint.)

(b) **Optional:** Let A, B, D be matrices and $\vec{b}, \vec{c}, \vec{d}, \vec{f}$ vectors (all of suitable dimensions). Convert the following LP with variables \vec{x} and \vec{y} (where \vec{x}, \vec{y} are vectors) into SEF:

$$\min\{\vec{c}^{\mathsf{T}}\vec{x} + \vec{d}^{\mathsf{T}}\vec{y}\}\tag{18}$$

subject to

$$A\vec{x} \ge \vec{b} \tag{19}$$

$$B\vec{x} + D\vec{y} = \vec{f} \tag{20}$$

$$\vec{x} \ge \vec{0}.\tag{21}$$

- 3. (Problem 1 in Sec. 2.3 of the text, p. 56.)
 - In this exercise, you are asked to repeat the argument in Section 2.3 with different examples.

(a) Consider the following LP:

$$\max\{(-1,0,0,2)\vec{x}\}\tag{22}$$

Subject to

$$\begin{pmatrix} -1 & 1 & 0 & 2 \\ 1 & 0 & 1 & -3 \end{pmatrix} \vec{x} = \begin{pmatrix} 2 \\ 3 \end{pmatrix}$$
(23)

$$\vec{x} \ge \vec{0}.\tag{24}$$

Observe that $\vec{x} = (0, 2, 3, 0)^{\mathsf{T}}$ is a feasible solution. Starting from \vec{x} , construct a feasible solution $\vec{x'}$ with value larger than that of \vec{x} by increasing as much as possible the value of exactly one of x_1 or x_4 (keeping the other variable unchanged).

(b) Consider the following LP:

$$\max\{(0, 0, 4, -6)\vec{x}\}\tag{25}$$

subject to

$$\begin{pmatrix} 1 & 0 & -1 & 1 \\ 0 & 1 & -3 & 2 \end{pmatrix} \vec{x} = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$$
(26)
$$\vec{x} > \vec{0}$$

$$\vec{x} \ge 0. \tag{27}$$

Observe that $\vec{x} = (2, 1, 0, 0)^{\mathsf{T}}$ is a feasible solution. Starting from \vec{x} , construct a feasible solution $\vec{x'}$ with value larger than that of \vec{x} by increasing as much as possible the value of exactly one of x_1 or x_4 (keeping the other variable unchanged).